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(54) **Network for transmitting isochronous-source data with a frame structure**

Netz zur Übertragung von isochronen Quelldaten mit einer Rahmenstruktur

Réseau de transmission de données source isochrones avec structure de trame

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Description

[0001] The present invention relates to a data communication network, such as a local area network or wide area network, and in particular to a network for transferring isochronous data.

BACKGROUND OF THE INVENTION

[0002] In general terms, isochronous data is data which is non-packetized and of indeterminate, potentially continuous duration. Increasing availability of multi-media computers and work stations has contributed to an increase in interest in the transfer of data from isochronous data sources and sinks. An isochronous data source is a device which outputs data in a continuous stream, usually at a substantially constant average data rate. Examples include video cameras, which output a substantially continuous stream of data representing images and associated sounds, and telephone output, which can be a substantially continuous output of voice data (either analog or digitized). An example of an isochronous data sink is a video monitor which can receive a substantially continuous stream of video data for display.

[0003] Related to isochronous sources and sinks is the concept of data transfer which can also be either isochronous or non-isochronous. One type of non-isochronous data transfer is a packet-type transfer. As shown in Fig. 1A, data can be transferred in a plurality of packets 12a, 12b which can be either constant-sized or variable-sized. Each packet includes a field of data 14a, 14b which may be preceded and/or followed by non-data information such as preamble information 16a, 16b housekeeping information such as data source information, data destination information, and the like 18a, 18b and a frame end marker 20a. As seen in Fig. 1A, because the fields provided for data 14a, 14b are not substantially continuous, the packetized scheme of Fig. 1A is not isochronous but is "bursty" in nature. An example of packetized data transfer is the commonly-used ethernet system, one implementation of which, known as 10BASE-T is described in the Draft Nine Supplement to IEEE Standard 802.3, dated November 15, 1989.

[0004] Another type of non-isochronous data transfer is a token ring system. In a token ring system, a node is permitted to transmit data only after receipt of an electronic "token." As depicted in Fig. 1B, a first station may transmit a token 22a which is received 24a by a second station whereupon the second station may begin transmission of data 26a. After a period of data transmission, the second station transmits the token 22b which is received by a third station 24b that can then begin its own transmission of data 26b. As seen in Fig. 1B, because data transmission is synchronized with the occurrence of an event (the arrival of a token), the token ring system is not an isochronous data transfer system. One commonly used token ring network is described in IEEE Standard 802.5.

[0005] In contrast, Fig. 1C schematically depicts isochronous data transfer. In isochronous data transfer, the data transfer or connection is initiated, such as by initiating a telephone conversation or beginning a video camera transmission 30. After the connection is initiated, transmission of the data, possibly accompanied by transmission of housekeeping information (such as destinations, audio or video timing and the like) is provided substantially continuously for an indeterminate period, such as until termination of the connection 32. Although it may be that not every bit transferred represents a data bit (since "housekeeping" bits may be also transferred), the transfer of data is substantially continuous in the sense that there are no substantial periods during which no data bits are transferred. It is possible that the data being transferred is "null" data such as silence during a telephone conversation or transfer of a blank video image. One type of isochronous data transfer is the Fiber Distributed Data Interface-II (FDDI-II) as described, for example, in FDDI-II Hybrid Multiplexer, Revision 2.4; dated March 25, 1991.

[0006] Previous attempts to accommodate isochronous data on a data network having resulted in characteristics which are disadvantageous for at least some applications. Some previous devices and methods have transmitted data using a frame structure. However, previous devices have not adequately transmitted data in continuous time frames while accommodating the data rate needs of a variety of data sources and sinks, including both isochronous and non-isochronous data sources and sinks. Previous devices have not provided for sufficient independence of various source and sinks. Previous devices have not adequately provided for guaranteed sustainable and switchable bandwidth for isochronous data, and particularly have not permitted a user to select a desired bandwidth for transmission of isochronous data. Previous devices have not adequately provided for interaction of network with an external clocked system and particularly have not adequately provided for contention between two or more different external clocks for accommodation to the data rate of a non-selected or skewed external clock reference.

[0007] Accordingly, it would be advantageous to provide an isochronous data communication system in which data is transmitted in a frame structure while accommodating differing data rate requirements of various data sources and sinks. It would also be advantageous to provide a system in which the bandwidth for given isochronous sources or sinks can be selected with a predetermined granularity. It would also be advantageous to provide a network which can accommodate its timing to an external clock reference that can co-ordinate output from the network of different systems running according to two different clocks or according to skewed clocks.

[0008] The state of the art is represented by the transactions of the Second IEEE National Conference on Telecom-

munications, April 1989 (York/GB), pages 173-178, which describes a multi-service local area network having an integrated voice and data interface. Similar subject-matter is disclosed in IEICE Transactions, Volume E 74, No. 9, September 1991 (Tokyo/JP), pages 2696-2702. The invention is defined in the claims, which are characterised with respect to the former of the aforementioned references.

[0009] The present invention provides for communication of data to and from isochronous data sources and sinks preferably in such a way that a predetermined bandwidth is available to an isochronous source/sink. Of the total bandwidth used for communication over the network links or physical media, at least a portion of the total bandwidth is dedicated to isochronous traffic. Preferably the bandwidth available for isochronous traffic can be selected or allocated with a predetermined granularity, e.g. so that the quality of transmission service desired for a given isochronous source or sink can be selected and the selected bandwidth can be sustained throughout the isochronous communication. Preferably a portion of bandwidth on the link is also dedicated to convey data to and from non-isochronous sources and sinks, as well as housekeeping information (such as information relating to data sources and destinations).

[0010] Preferably, the system of the present invention can be provided so that it is transparent to previously-available Media Access Controllers (MACs) and applications.

[0011] In one embodiment, data from an isochronous data source is time-division multiplexed with the data output from at least a second source using a predetermined frame structure. The second data source can include non-isochronous-sourced data, maintenance "M channel" data and connection control "D channel" data such as destination, source bandwidth and/or status information. Preferably, all four types of data are accommodated in the frame structure. The frame structure or template is constructed to satisfy the data rate requirements of the various data sources and sinks. The various types of data are transmitted in a substantially independent manner and, in particular, the occurrence of, e.g., data collision (of ethernet data) or a token loss (in token-ring protocol data) will not affect transmission of the isochronous data which can continue to be sent at an unchanged data rate despite the occurrence of data interruptions (e.g., collision or token loss) in other data types.

[0012] Because each frame has a pre-determined number of time slots for each type of data, and because the time frames repeat at a pre-determined interval, the effective data rate for any one of the types of data carried is constrained. According to one embodiment of the present invention, the frame has one or more time slots which can be used for adjusting the data rate to the particular data rate requirements of various data sources and sinks. For example, one or more time slots can be designated as rate-adjustment slots. The rate-adjustment slots can be used to carry the specified type of data during some time frames but do not carry that specified type of data during other time frames. For example, by alternating a rate adjustment time slot between "data" and "no data" a data rate adjustment which is equal to one half of the bandwidth represented by a single time slot can be achieved.

[0013] Preferably, the present system is implemented as a star-topology network with data sources transmitting to a central hub which, in turn, transmits the data to data sinks. A single node can act as both a source and a sink. Several such star-topology systems can be connected by providing interconnection of the hubs, for example, in a ring structure (Fig. 2) or to provide a tree structure. The multiplexed data which arrives at the hub is de-multiplexed to separate the isochronous-source data, the non-isochronous-source data and the D channel and maintenance information. The non-isochronous-source data can be provided to hub circuitry specialized for handling the non-isochronous data stream. The separated isochronous data is conveyed to circuitry configured for handling the isochronous data and transmitting to the destination nodes of the network or to connected hubs. Preferably, the hub has sufficient intelligence to set up and maintain isochronous communication sessions or calls which may be requested on the D channel.

[0014] The hub contains multiplexers for combining both isochronous-sourced data and other data, e.g. non-isochronous-sourced data. These data sources, along with M channel and D channel data, are multiplexed in a fashion similar to the multiplexing which occurred at the nodes and the multiplexed data is transmitted back to the nodes, preferably over a separate set of one-way twisted pair media. The nodes contain de-multiplexers, similar to those found in the hub.

[0015] The frame structure used for transmission over the media provides a data rate which, while substantially constant over a long time frame, is variable over a short time period (such as less than one template) and buffering can be used for smoothing the variable data rate to provide compatibility with the data rates of the various sources and sinks.

[0016] In one embodiment, the timing of the frame transmissions can be coordinated with a reference clock signal which is external to the network, (e.g., a reference clock for a different network or system.) For example, the external reference clock could be the reference clock of a public telephone system or other wide area network ("WAN"). The external reference clock can be provided to the network at the hub or at any node. The clock signal is then propagated throughout the rest of the network. When more than one external reference clock is available (e.g., when a local area network according to the present invention is coupled to two or more wide area networks) one of the external reference clocks is selected as the reference clock (e.g., using a contention system such as on FDDI, by station management (SMT)). In the present invention, it is possible that a given node may transmit at a frame rate which differs slightly from a propagated reference clock. In this case, the frame length may be shortened or lengthened as needed for adjustment.

In cases where two or more separately-clocked systems are connected to the network (e.g., two or more WAN's), a buffer can be provided for re-timing one or more of the output.

BRIEF DESCRIPTION OF THE DRAWINGS

[0017]

Figs. 1A, 1B and 1C of the timing of a packet transmission system, a token ring transmission system, and an isochronous transmission system respectively.

Fig. 2 is a schematic block diagram showing three nodes connecting to a hub card according to one embodiment of the present invention;

Fig. 3 is a schematic block diagram showing a number of hubs connected together using a ring structure;

Fig. 4 is a schematic block diagram showing the relationship of timing signals from two wide area networks with a network according to the present invention;

Fig. 5 is a partial block diagram depicting the relationship of a node FIFO and a hub FIFO to data sources, sinks and repeaters, according to an embodiment of the present invention;

Fig. 6 is a block diagram depicting the packet receive interface, according to an embodiment of the present invention;

Fig. 7 is a schematic block diagram of node circuitry for multiplexing and preparing data for transmission over the media and for receiving information from the media and demultiplexing the data; and

Fig. 8 is a timing diagram showing the relative timing of transmissions and receptions at the hub and nodes.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENT

[0018] According to an embodiment of the invention, a data communication system for isochronous data can be configured in a star-topology with a plurality of nodes 42a, 42b, 42c (Fig. 2), each coupled to a hub 44a by data links comprising physical data transmission media such as one-way twisted pair wires 46a-46f. The number of nodes can be adjusted depending on the data transmission needs and objectives of the system. In one embodiment, each hub is configured to accommodate connection with up to 16 nodes. Each node 42a, 42b, 42c includes circuitry 50a, 50b, 50c for receiving data, converting it to a form suitable for transmission onto the physical media 46a, 46c, 46e and receipt of signals from the physical media 46b, 46d, 46f and conversion to a form suitable for use by the data sinks. Each of the nodes 42a, 42b, 42c includes data sources and sinks 48a-48g. The data sources and sinks can be isochronous sources and sinks such as video cameras 48a, 48d and monitors 48b, 48e, non-isochronous sources and sinks such as an ethernet media access controller 48c, 48g, and control signaling or D channel sources and sinks such as an emulated or virtual key pad 48f provided, for example, on a personal computer (PC) terminal. Each of the nodes 42a, 42b, 42c can include various types of sources and sinks, e.g., strictly isochronous sources and sinks, such as depicted for node one 42a, strictly non-isochronous sources/sinks as depicted for node three 42c or both isochronous and non-isochronous sources and sinks as depicted for node two 42b. The physical layer 52 of the network system depicted in Fig. 2 includes the node data receivers and converters 50a, 50b, 50c, the physical media 46a-46f and the hub 44a.

[0019] The hub 44a includes circuitry 54a, 54b, 54c for receiving data from the physical media 46a, 46c, 46e separating the isochronous-sourced data from the non-isochronous-sourced data and the D channel and M channel data and converting separated data into a form suitable for handling by downstream hub circuitry 56. In the depicted embodiment the separated isochronous-sourced data is provided to a time slot interchange controller for placing the data on a high-bandwidth bus so that it can be transported to and retrieved by hub circuitry 54a, 54b, 54c for transmission to various destination nodes 42a, 42b, 42c. The separated non-isochronous data is provided to circuitry 60 configured to convey the non-isochronous data to the hub circuitry 54a, 54b, 54c for transmission to destination nodes 42a, 42b, 42c. In an embodiment in which non-isochronous-sourced data includes ethernet data, the hub circuitry 60 can be a standard ethernet repeater processor. In this way, a system which incorporates the present invention can be at least partially backwards-compatible with previous ethernet hub systems. The D channel and M channel data is provided to a signaling processor 62 which performs various maintenance and control functions such as identifying and alerting users of error conditions, and setting up requested calls, i.e. source/destination paths e.g. by communicating with the isochronous and non-isochronous controllers 58, 60, e.g. over data path 64.

[0020] Data sent from an isochronous device, e.g., 48d is a continuous stream of digitized data from such as a video camera. In one example, the data from isochronous device 48d will be taken as having a data rate equal to the American "T1" Standard of 1.544Mbit/sec. Data output from the ethernet MAC 48c is provided at the standard 10BASE-T ethernet rate of 10Mb/sec. D channel data is provided from the host processor or the ethernet MAC 48c or, for example, from the virtual key pad 48f at a variable data rate, such as a rate not exceeding about 64Kb/sec. These incoming data

streams are provided over lines 66a, 66b, 66c to node circuitry 50b. The incoming data from the various sources is provided to a multiplexer 70 which performs time-division multiplexing on a four-bit basis. The pattern for the time-division multiplexing is a repeating series of frames or templates. Preferably, the frames are repeated with a period which is the same as the public telephone network, (i.e., 125 microseconds). A clock signal is used for timing the frame transmissions, i.e., for establishing the points, (i.e., 125 microseconds), at which the transmission of each frame begins. According to one embodiment, the reference clock is initiated at one physical location in the network and propagated throughout the network. When the network is connected to an external system such as a public telephone network or other wide area network, the clock signal from the external system can be used to establish the reference clock for the network of the present invention. Thus, in one embodiment a hub or node is connected to a public telephone network and the frame sync from the public telephone network is propagated through the network of the present invention.

[0021] In some instances, as depicted in Fig. 4, the network of the present invention 502 may be connected to more than one wide area network, (e.g., by connecting a first node 504 to a wide area network 506, and a second node 508 to a second wide area network 510.) In this case, one of the connections, e.g., 506 can be chosen, e.g., by the system administrator, as the reference of the network of the present invention. Communication with the other (non-selected) wide area network 510 is conducted by a buffer or a data storage device 512 containing multiple frames of data. The buffer 512, is preferably large enough to contain multiple frames of data in order to accommodate the phase and/or frequency skews between the frame structure of the network of the present invention 502, and the frame structure of the external system 510.

[0022] Table I depicts the manner in which the various data streams, and additional data and control bytes are time-division multiplexed. In the embodiment depicted in Table I, each frame is divided into four blocks with each block consisting of eight groups. There are 16 signals in each group, thus providing 512 paired symbols in each frame. Each symbol in Table I represents four bits of data so that every group of two symbols represents one 8-bit byte of data. In Table I, E represents four bits of data from the ethernet stream 16b, B designates four bits of data from the isochronous stream 66a, D represents four bits of data from the signaling or D channel stream 66c; and M represents four bits of maintenance data as described below. In addition, certain byte-length patterns are provided. JK represents a frame synchronization pattern and EM (the first two bytes of block three in Table I) represents an ethernet "pad" followed by a maintenance nibble as described below. Thus, the D channel time slots are found in the 257th and 248th time slots. The maintenance time slots are found in the 129th, 130th and 386th time slots. By positioning the D channel, M channel and frame synchronization symbols at the beginning of each block, the present invention is able to achieve a structured frame template usable with various non-isochronous protocols, with only minor modifications required. The resultant symmetry eases generation of the template. As seen in Table I, each frame contains 256 bytes which can be considered in thirty-two groups of eight bytes each, or four blocks of sixty-four bytes each. As seen in Table I, each block is headed by a symbol pair. The JK symbol pair is a unique synchronization pattern which indicates the beginning of the frame. The second block's header is the maintenance byte. The maintenance byte is used to convey information such as notification of the existence of a transmission error, or other signals to be conveyed to the far end of the link, such as other necessary physical layer information. The third block's header is the location of the D channel. The D channel is provided as a service to the user. The software and hardware can be configured to use the bandwidth of the D channel for a variety of purposes. It is anticipated that the D channel will be used for connection set-up and tear-down, which includes transmitting information related to the source node address and intended destination node address of data. Thus, the D channel can contain information which may be used by the hub to set-up calls, (i.e., to route information from a source to a destination.) The symbol pair heading the final block has a first symbol which is reserved for a networking scheme symbol. This symbol can be used as the "rate adjustment" time slot. Thus, it can be used either to provide data or "housekeeping" information in this time frame (termed an "over frame") or can carry a symbol indicating that no data is being transmitted in this time slot for this time frame (termed an "under frame"). The types of symbols which may be used for "housekeeping" information are specific to the type of data source or sink. For example, if the data source or sync is a token ring system, the housekeeping symbol might be the frame start delimiter. If the data source or sync is an ethernet system, the housekeeping symbol might represent "no-carrier" or "un-aligned data". Some networking schemes use the absence of carrier as an indication of the link state. In the ethernet protocol, the absence of carrier issues to mark the boundaries between packets. In the described frame structure, however, there is never an absence of carrier. Transmission is substantially continuous. Therefore, a "no carrier" state can be conveyed across the link using a unique symbol pattern outside the O-F set of data patterns. These possibilities are described in greater detail in EP-A-0 596 645.

TABLE I

BLOCK 0:															
J	K	E	B	E	B	E	B	E	B	E	B	E	B	E	Group0

TABLE I (continued)

BLOCK 0:																
5	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group1
	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group2
	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group3
	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group4
	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group5
	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group6
	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group7
BLOCK 1:																
15	M	M	E	B	E	B	E	B	E	B	E	B	E	E	E	Group8
	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group9
	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group10
	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group11
	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group12
	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group13
	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group14
20	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group15
BLOCK 2:																
25	D	D	E	B	E	B	E	B	E	B	E	B	E	E	E	Group16
	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group17
	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group18
	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group19
	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group20
	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group21
	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group22
30	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group23
BLOCK 3:																
35	E	M	E	B	E	B	E	B	E	B	E	B	E	E	E	Group24
	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group25
	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group26
	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group27
	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group28
	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group29
	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group30
40	E	E	E	B	E	B	E	B	E	B	E	B	E	E	E	Group31
JK = Frame Synchronization Pattern MM = 8 Maintenance Bits D = D Channel EM = Ethernet Pad & 4 Maintenance Bits E = Ethernet Packet Channel B = Isochronous Channel																
45																
50																

[0023] If the rate adjustment frame is an "over frame", (i.e., one in which data or housekeeping information is conveyed) then the symbol is one of the following: (a) housekeeping symbol; or (2) the data symbol representing one of the hexadecimal values 0 through F. If the symbol is an "underframe", then a "no data symbol" is transmitted. In one embodiment, "no data" symbols are alternated with the other possibilities on an every-other-frame basis to provide an average of two bits of sourced information per frame cycle for this time slot. It would also be possible to provide other combinations such as transmitting "no data" one out of three frames, two out of three frames, one out of four frames, three out of four frames, and so forth. This allows the channel rate for a particular data to be "adjusted" precisely to

the required amount for the networking scheme. Where more than two types of data sources each require "adjustment", (i.e., where a network carries both ethernet data and token ring data), one or more rate adjustment time slots can be provided for each of these data sources and sinks.

[0024] The second symbol in the symbol pair heading the fourth block consists of four additional maintenance bits. Together with the 8 maintenance bits from the second block, there are 12 maintenance bits which are used for physical layer signaling.

[0025] In one embodiment, each node or transmitter contains its own local crystal oscillator for determining the lengths of the time slots and, thus, the frames. Preferably, the data communication system as a whole is timed according to a reference clock signal which can be propagated through the network. Because of small variations in accuracy of the local crystal oscillators and/or the reference clock, it may happen that there are small frequency differences between the local crystal oscillator and the network reference clock signal such that, for example, the frames at a given node or transmitter are somewhat longer or shorter than the ideal frame length determined from the reference clock signal. In such a situation, according to one embodiment of the invention, the difference between the actual frame length (determined by the local crystal oscillator) and the desired or ideal frame length (determined by the network reference clock signal) is detected, and an adjustment is made to the frame length as often as needed to effectively shorten or lengthen the frame. In this embodiment, if the local clock is slow compared to the network clock reference, the last symbol of a frame is dropped (i.e., the last symbol of Group 31 depicted in Table I) to provide a frame having only 511 symbols. The data which would have been sent in the 512th symbol in a normal frame will then be added in the next sequential unused rate adjustment or "pad" time slot. If the local oscillator is fast, compared to the network reference clock, an additional symbol is added to a frame to produce a frame having 513 symbols. The added symbols is preferably added at the end of the normal 512 symbols of a normal frame. By detecting variances between the local clock and the reference clock and shortening or lengthening the frame as needed for adjustment, the present invention permits transmission of frames while relying on local oscillators which are substantially consistent with frame transmission rates determined by a network clock reference, without having to provide circuitry for strictly synchronizing a node or local transmitter to the network reference clock (such as by using a phase-lock loop system).

[0026] The structure imposed by the frame or template introduces an amount of jitter. In particular, when non-isochronous-sourced data, such as packet data is repositioned in time so as to fill the proper time slot for transmission, according to the frame structure of Table I, jitter is introduced. As depicted in Fig. 5, if data is to be repeated in the hub 44 upon reception 60, a small first-in-out (FIFO) memory 330 can be used for re-timing the data to eliminate the effects of the introduced jitter and provide the timing expected by the repeater circuitry 60. At the transmitting end, if the data to be transmitted onto the physical medium 46 is arriving from a source 48, whose rate cannot be directly controlled, the transmitted data can also be passed through a FIFO 332.

[0027] Preferably, the frame structure is provided so as to prevent an unacceptably large passage of time between sequential transmissions of particular data types in a given time frame. In the frame structure of Table I, there are no more than two time slots between any two sequential non-isochronous-sourced time slots in a single time frame and there are not more than five time slots between any two sequential isochronous-sourced time slots in a single time frame.

[0028] The depicted frame structure provides an allocation of bandwidth such that the data rate for the isochronous and non-isochronous data are compatible with, e.g., 10BASE-T ethernet data rates. Other types of frame structures could be used in connection with other isochronous and/or non-isochronous non-isochronous data sources and sinks such as token ring sources and sinks, in which case a different frame structure can be used to provide an allocation of bandwidth suited for the particular purpose.

[0029] The time-multiplexed data is then encoded by an encoder 72. In the depicted embodiment, the encoder performs four/five encoding. One particular form of four/five encoding conforming partially to the ANSI X3T9.5 standard, is depicted in Table II. The five bit code symbols of the particular four/five implementation have been chosen so as to maintain the AC balance of the physical medium, and to minimize the frequency spectrum of wave forms as they are transmitted along the physical medium. This eases the task of the data decoder, which is typically a phase lock loop device, in recovering the data and the transmission clock. These patterns, when properly combined, have a maximum of three bit times with no transition. With a bit rate of 48.8 ns, the minimum transition rate is 3.41 MHz or 146.5 ns between transitions.

TABLE II

Symbol	Encoded (5 bit)	Decoded (4 bit)	Description
0	11110	0000	Data 0
1	01001	0001	Data 1
2	10100	0010	Data 2

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TABLE II (continued)

Symbol	Encoded (5 bit)	Decoded (4 bit)	Description
3	10101	0011	Data 3
4	01010	0100	Data 4
5	01011	0101	Data 5
6	01110	0110	Data 6
7	01111	0111	Data 7
8	10010	1000	Data 8
9	10011	1001	Data 9
A	10110	1010	Data A
B	10111	1011	Data B
C	11010	1100	Data C
D	11011	1101	Data D
E	11100	1110	Data E
F	11101	1111	Data F
I	11111	1010	No Ethernet Carrier
S	11001	0111	No Ethernet Data
V	01100	0010	Unaligned Data
T	01101	0101	Unassigned
J	11000	1101	Frame Sync Part 1
K	10001	1101	Frame Sync Part 2
Q	00000	0010	Invalid
H	00100	0001	Invalid
R	00111	0110	Invalid
V	00001	0110	Invalid
V	00010	0010	Invalid
V	00011	0010	Invalid
V	00101	0010	Invalid
V	00110	0010	Invalid
V	01000	0010	Invalid
V	10000	0010	Invalid

[0030] Comparing the standard ethernet signalling rate and data rate to that provided by the frame structure of Table I and encoding of Table II, it is seen that in the present scheme, a frame of data contains 313 "E" symbols or 1252 E bits. By transmitting frames at a rate of one frame every 125 microseconds, the present scheme has a capacity for transmitting 10.016 Mbits/sec of ethernet-sourced data interspersed with 6144 Kbits/sec of isochronous data.

[0031] As one example, it is assumed that the data rate output from the isochronous source 48d is 1.544 Mbits/sec. However, the frame structure noted above provides an isochronous bandwidth capability of 6.144 Mb/s. Thus, the single isochronous source 48b in the present example can be entirely accommodated using only 48 of the 192 "B" symbols per frame (i.e. 24 bytes per frame or 192 bits per frame). This leaves 72 "B" symbols per frame unutilized by source 48d. Thus, it would be possible to use the isochronous channel of this embodiment of the invention to carry isochronous data from four isochronous sources, each outputting data at a rate of 1.544 Mb/sec. It would also be possible to use the described embodiment to carry the output from three isochronous sources, each outputting data at a rate of 2.048 Mb/sec (i.e. the CEPT European standard). A basic rate ISDN channel could be supported by using three 64 Kb/s slots within the isochronous channel or by e.g., 2 B slots and the D channel. It would also be possible to transmit data from up to 96 isochronous sources, each outputting data at a rate of 64 Kb/s. Thus, it is possible to use the present invention to transfer data from a single isochronous source outputting at a rate of 6.144 Mb/s or the isochronous bandwidth available may be allocated with a granularity of 64 Kb/s, i.e., it may be split into multiples of 64 Kb/s.

[0032] The data output is sent to a transmitter or driver 78a and the signal is transmitted over the physical medium 46c. The physical medium 46c can be any of a number of media types including twisted pair, coaxial or fiber optic cable.

[0033] The data sent over the physical media 46a is received in the hub 44a. The hub contains a plurality of circuit devices 54a, 54b, 54c, each one coupled to one of the nodes 42a, 42b, 42c by the physical media 46. Link detect

circuitry receives the data from the physical media 46 for detection of the mode in which the node is operating, (e.g., 10BASE-T or isochronous ethernet) and outputting a mode select signal, as described more fully in EP-A-0 596 648. A demultiplexer 92 separates the received data into the isochronous-sourced data 94a the non-isochronous-sourced data 94b and signaling data, such as D channel and maintenance data 94c.

[0034] Both the non-isochronous-sourced data 94b and the isochronous-sourced data 94a are made available to the various hub circuitry components 54a, 54b, 54c, as needed for transmission back to destination nodes.

[0035] Circuitry 58 can also be provided with a local loop-back capability. The local loop-back capability permits circuitry 58 to transfer data directly from the receive buffer 138a to the transmit buffer 154 without first placing the data onto the TSI ring 134, thus freeing TSI bus bandwidth. The multiplexer 156 can be used to control loop-back. Local loop-back capability is described more thoroughly in EP-A-0 596 452.

[0036] After the hub has received data from a node and conveyed it to another node for transmission to the destination node, the data which is intended for the destination node must be placed in a form suitable for transmission across the physical media 46. The data sent from the hub 44a to the nodes 42 is sent in a frame format which is preferably substantially the same as the frame format used for the data sent from the nodes 48 to the hub 44a as described above.

At the nodes 42, the circuitry 50 includes devices (Fig. 7) for decoding and demultiplexing data. Decoded and demultiplexed data is then delivered to the various data sinks in the nodes 42.

[0037] In one embodiment, it is desirable to control the timing of the transmissions from the nodes 48 to the hub 44 and those from the hub 44 to the nodes 48 to assist in reducing delay and minimizing the amount of buffering memory required. In particular, it is desired that the hub 44 should be able to transmit to the nodes, over media 46a, 46c, 46e at about the same time that data transmitted from the nodes over media 46b, 46d, 46f is received at the hub. Fig. 8 depicts a scheme for achieving this type of timing. As shown in Fig. 8, the timing can be synchronized with a 125 microsecond reference clock signal 214. In this example, the reference signal 214 provides an ascending clock edge every 125 microseconds. The reference signal can be provided by any of a number of sources. Preferably, an embodiment of the present invention is configured to permit a reference signal 214 to be synchronized to an external clock reference, such as a reference signal from a wide band network or from a FDDI-II ring.

[0038] Each node and the hub device will contain its own high-accuracy transmission encoding clock, typically a crystal oscillator. Thus, when viewed at the bit level, the links 46 are synchronous to each other, being referenced only to the sourcing clock oscillator. Thus, according to an embodiment of the present invention, the synchronous frame structure is conveyed over synchronous links.

[0039] At the beginning of a cycle, the hub 44 will begin to transmit a frame to the node, as indicated by the timing marks on time line 216. Because of the line delays in the physical media, the time at which the nodes will receive the frame transmitted by the hub will lag the time when they were sent out from the hub as shown by time line 218.

[0040] The timing scheme described guarantees that the cycle received from the node will arrive slightly sooner than the next cycle will be transmitted from the hub. A small FIFO can be inserted into the hub's received data stream to accurately align the arrival of the cycle. In similar applications, it will also be advantageous to provide a FIFO in the node which would hold the data synchronized with the received cycle reference until it is to be transmitted. Provision of these FIFOs is described in more detail in EP-A 0 596 650.

[0041] In light of the above description, a number of advantages of the present invention can be seen. The present invention is able to provide a frame-structured data transfer which can simultaneously accommodate the data rate needs of several types of data sources and sinks, preferably both isochronous sources and sinks and non-isochronous sources and sinks. Preferably, the frame structure also provides a bandwidth for maintenance information, D channel information, and frame synchronization information. Rate adjustment time slots in the data frames allow the data rates to be adjusted to the rates particular to various sources and sinks. Jitter introduced by the frame structure can be buffered using FIFO memories between the data links and the sources and sinks or hub repeaters. The timing of frame transmissions can be coordinated with clocks in external systems such as public telephone systems or other wide area networks. The frame structure permits one, several or all time slots intended for isochronous-sourced data to be allocated to a given isochronous source or sink. This permits bandwidth select ability and maintenance ability for the isochronous data, preferably with a pre-determined granularity such as 64Kb granularity.

[0042] A number of modifications and variations of the invention can also be used. Although the frame structure of Table I is believed to be particularly useful for the mix of data anticipated in normal office or business installations having both isochronous and ethernet sources and sinks, other frame structures can also be provided. For example, in a situation where greater bandwidth is desired for ethernet traffic, it would be possible to provide a larger number of "E" time slots and a fewer number of one or all of the "B", "D", or "M" time slots. Further, if a token ring station was provided as one or more of the non-isochronous data sources or sinks, a number of non-isochronous-sourced time slots would be provided to accommodate the data rate needs of a token ring system. For example a isochronous-token ring system might require a larger bandwidth, such as a 22.4 Mbit/sec link. Such a system could be provided with a frame size of about 700 symbols so as to not exceed the 32 Mbit/sec signaling rate of typical token ring media. A frame structure could be provided with longer or shorter duration frames and/or longer or shorter duration time slots although

the disclosed frame and time slot durations are believed particularly useful because the mapping non-isochronous and isochronous bandwidths maps into existing networking schemes, such as ethernet, ISDN, ATM and/or SONET. The disclosed frame and time slot durations provide the 10 Mbit/s bandwidth needed, e.g., for ethernet and providing a multiple of twelve 8 KHz channels, e.g., for isochronous, plus bandwidth, e.g., for controlling the link and isochronous connections. Although the disclosed frame structure provides for 64 Kbit/s of granularity for isochronous data, other granularities could be provided by altering the duration of time slots, the duration of time frames, or the number of data bits transmitted per time slot. The timing for frame transmissions could be derived from sources other public telephone network or other wide area networks, such as an internal reference clock T1 (1.544 Mbit/s) link, or ISDN connection, or asynchronous transfer mode (ATM). It is possible to use some aspects of the invention without using other aspects. For example, it would be possible to use the frame structure of Table I without providing a buffer for accommodating a second external clock.

Claims

1. Apparatus for transmitting information onto a physical medium (46) of a local area network comprising: means (50) for receiving a plurality of inputs, including at least a first isochronous data stream and a second stream of data; and means (78) for time-multiplexing information from said plurality of inputs by transmitting times frames in a repeating pattern substantially without pause between said time frames, each time frame having a plurality of time slots, the time frames having substantially equal durations, and each time slot being used to transmit a predetermined number of data bits; wherein first predetermined ones of said time slots of each time frame are used for transmitting data bits from said first isochronous data stream and second predetermined ones of said times lots of each time frame are used for transmitting data bits from said second stream of data; characterised by means for detecting a difference between an actual frame length set by a local oscillator and a frame length determined by a network reference clock signal and shortening or lengthening a frame by removing or adding a time slot according as the local oscillator is slow or fast relative to the reference clock.
2. Apparatus according to claim 1, wherein said second stream of data is a non-isochronous-sourced data stream, said first isochronous data stream is transmitted at a first data rate and wherein said first data rate is independent of interruptions in transfer of said second stream of data.
3. Apparatus according to claim 1, wherein said second stream of data is ethernet-sourced data, having a potential for data-collision events, and said first isochronous data stream is transmitted at a data rate which is substantially unchanged in response to the occurrence of a data-collision event in said second stream of data.
4. Apparatus according to claim 1, wherein said second stream of data is token-ring-sourced data, having a potential for a token-loss event, and said first isochronous data stream is transmitted at a data rate which is substantially unchanged in response to the occurrence of a token-loss event in said second stream of data.
5. Apparatus according to claim 1, wherein said second stream of data is a non-isochronous-sourced data stream and wherein at least one of said time slots is used for transmitting data from said non-isochronous data stream during at least a first plurality of time frames, and said one of said time slots is used to transmit a non-data signal during a second plurality of time frames.
6. Apparatus according to claim 5, wherein said means for time-multiplexing transmits a predetermined pattern of said first plurality of frames and said second plurality of frames.
7. Apparatus according to claim 6, wherein said predetermined pattern is an alternation between said first plurality of frames and said second plurality of frames.
8. Apparatus according to claim 1, wherein said second stream of data is a non-isochronous-sourced data stream and wherein there are no more than two time slots between any two sequential non-isochronous-sourced time slots in a single time frame.
9. Apparatus according to claim 1, wherein there are no more than five time slots between any two sequential isochronous-source-filled time slots in a single time frame.
10. Apparatus according to claim 1, wherein the first and second time slots in each of said plurality of frames is used

for transmitting frame synchronization symbols.

11. Apparatus according to any foregoing claim and including a plurality of nodes, each connected to a hub, and further comprising means for receiving a reference clock signal, means for propagating said reference clock signal to said hub and to each of said nodes, and means for initiating transmission of said time frame at a predetermined time in relation to said propagated reference clock signal.

12. Apparatus according to claim 11 and including means for receiving a second reference clock signal, and means for selecting said first clock signal as a signal for use by said means for initiating transmission.

13. Apparatus according to claim 12, further comprising a buffer for storing data transmitted on said physical medium to permit output from said network synchronized with said second reference clock signal.

14. A method for transmitting information onto physical media coupled to a plurality of data source/sinks of a local area network, including at least a first isochronous data source/sink, the method comprising:

receiving, at a first of said plurality of data source/sinks, a plurality of inputs, including at least a first isochronous data stream and a second stream of data, said first isochronous data stream including at least data from said first isochronous data source/sink;

time-multiplexing information from said plurality of inputs by transmitting said information onto said physical medium in a repeating pattern of time frames substantially without pause between time frames, each time frame having a plurality of time slots, each time frame having a substantially equal duration, and each time slot being used to transmit a predetermined number of data bits;

a first plurality of said time slots being reserved for transmitting only data bits from said first isochronous data stream and second predetermined ones of said time slots being used for transmitting data bits from said second stream of data;

and allocating data from said first isochronous data source/sink to predetermined ones of said first plurality of time slots to provide a predetermined data bandwidth for isochronous data from said first isochronous data source/sink;

characterised by detecting a difference between an actual frame length set by a local oscillator and a frame length determined by a network reference clock signal and shortening or lengthening a frame by removing or adding a time slot according as the local oscillator is slow or fast relative to the reference clock.

Patentansprüche

1. Vorrichtung zur Übertragung von Information über ein physikalisches Medium (46) eines lokalen Netzwerks mit:

Mitteln (50) zum Empfang einer Mehrzahl Eingaben, die zumindest einen ersten isochronen Datenstrom und einen zweiten Datenstrom aufweisen, und

Mittel (78), um die Information der Mehrzahl Eingaben mit Zeitrahmen in einem sich wiederholenden Muster in der Zeit zu multiplexen, wobei im wesentlichen keine Pause zwischen den Zeitrahmen liegt, jeder Zeitrahmen eine Mehrzahl Zeitfenster aufweist, die Zeitrahmen im wesentlichen gleiche Dauer haben und jedes der Zeitfenster zur Übertragung einer vorgegebenen Zahl von Datenbits verwendet wird, wobei erste vorgegebene Zeitfenster jedes Zeitrahmens zur Übertragung von Datenbits des ersten isochronen Datenstroms verwendet werden, und zweite vorgegebene Zeitfenster jedes Zeitrahmens zur Übertragung von Datenbits des zweiten Datenstroms verwendet werden;

gekennzeichnet durch

Mittel zur Detektion einer Differenz zwischen einer tatsächlichen Rahmenlänge, die durch einen lokalen Oszillator eingestellt wird, und einer Rahmenlänge, die durch ein Netzwerkbezugstaktsignal bestimmt wird, und zum Verkürzen oder Verlängern eines Rahmens durch Entfernen oder Hinzufügen eines Zeit fensters in Abhängigkeit davon, ob der lokale Oszillator relativ zum Bezugstakt langsam oder schnell ist.

2. Vorrichtung nach Anspruch 1, bei der der zweite Datenstrom ein Datenstrom einer nicht isochronen Quelle ist, der erste isochrone Datenstrom mit einer ersten Datenrate übertragen wird und die erste Datenrate unabhängig von Unterbrechungen bei der Übertragung des zweiten Datenstroms ist.
- 5 3. Vorrichtung nach Anspruch 1, bei der der zweite Datenstrom von einer Ethernet-Quelle mit der Möglichkeit von Datenkollisionsereignissen stammt und der erste isochrone Datenstrom mit einer Datenrate übertragen wird, die im wesentlichen unverändert bleibt, in Abhängigkeit von dem Auftreten eines Datenkollisionsereignisses in dem zweiten Datenstrom.
- 10 4. Vorrichtung nach Anspruch 1, bei der der zweite Datenstrom von einer Belegungsbitdaten-Quelle mit der Möglichkeit eines Belegungsbit-Verlustereignisses stammt und der erste isochrone Datenstrom mit einer Datenrate übertragen wird, die im wesentlichen unverändert bleibt, in Abhängigkeit von dem Auftreten eines Belegungsbit-Verlustereignisses in dem zweiten Datenstrom.
- 15 5. Vorrichtung nach Anspruch 1, bei der der zweite Datenstrom ein Datenstrom einer nicht isochronen Datenquelle ist und zumindest eines der Zeitfenster verwendet wird, um Daten von dem nicht isochronen Datenstrom während zumindest einer Anzahl von Zeitrahmen zu übertragen, und wobei dieses eine Zeitfenster verwendet wird, um ein Nicht-Datensignal während einer zweiten Anzahl Zeitrahmen zu übertragen.
- 20 6. Vorrichtung nach Anspruch 5, bei der das Mittel für das Zeitmultiplexen ein vorgegebenes Muster der ersten Anzahl von Rahmen und der zweiten Anzahl von Rahmen überträgt.
7. Vorrichtung nach Anspruch 6, bei der das vorgegebene Muster ein Wechsel zwischen der ersten Anzahl von Rahmen und der zweiten Anzahl von Rahmen ist.
- 25 8. Vorrichtung nach Anspruch 1, bei der der zweite Datenstrom ein Datenstrom einer nicht isochronen Datenquelle ist und nicht mehr als zwei Zeitfenster zwischen je zwei aufeinanderfolgenden Zeit fenstern der nicht isochronen Quelle in einem einzelnen Zeitrahmen liegen.
- 30 9. Vorrichtung nach Anspruch 1, bei der nicht mehr als fünf Zeitfenster zwischen zwei aufeinanderfolgenden, durch die isochrone Quelle gefüllten Zeitfenstern in einem einzelnen Zeitrahmen liegen.
10. Vorrichtung nach Anspruch 1, bei der das erste und zweite Zeitfenster in jedem der Mehrzahl von Rahmen zur Übertragung von Rahmen-Synchronisationssymbolen verwendet werden.
- 35 11. Vorrichtung nach einem der vorangehenden Ansprüche und mit einer Mehrzahl Knoten, die jeweils mit einer Nabe verbunden sind, und die desweiteren Mittel zum Empfang eines Bezugsaktsignals, Mittel zur Übertragung des Bezugsaktsignals an die Nabe und an jeden der Knoten, und Mittel zur Initiierung der Übertragung des Zeitrahmens zu einer vorgegebenen Zeit in bezug auf das ausgebreitete Bezugsaktsignal aufweist.
- 40 12. Vorrichtung nach Anspruch 11, die Mittel zum Empfang eines zweiten Bezugsaktsignals und Mittel zur Auswahl des ersten Taktsignals als Signal zur Verwendung durch das Übertragungsinitiierungsmittel aufweist.
- 45 13. Vorrichtung nach Anspruch 12, die desweiteren einen Puffer zum Speichern der über das physikalische Medium übertragenen Daten aufweist, um zu ermöglichen, daß die Ausgabe des Netzwerks mit dem zweiten Bezugsaktsignal synchronisiert ist.
- 50 14. Verfahren zur Übertragung von Information über ein physikalisches Medium, das mit einer Mehrzahl Datenquellen/Senken eines lokalen Netzwerks verbunden ist, die zumindest eine erste isochrone Datenquelle/Senke aufweisen, wobei das Verfahren umfaßt:

Empfang an einer ersten Mehrzahl Datenquellen/Senken einer Mehrzahl von Eingaben, die zumindest einen ersten isochronen Datenstrom und einen zweiten Datenstrom umfassen, wobei der erste isochrone Datenstrom zumindest Daten der ersten isochronen Datenquelle/Senke aufweist;

Zeitmultiplexen der Information der Mehrzahl Eingaben zur Übertragung der Information über das physikalische Medium in einem sich wiederholenden Muster aus Zeitrahmen, im wesentlichen ohne Pause zwischen den Zeitrahmen, wobei jeder der Zeitrahmen eine Mehrzahl von Zeitfenstern hat, jedes Zeitfenster im wesentlichen die gleiche Länge hat und jedes Zeitfenster zur Übertragung einer vorgegebenen Anzahl von Datenbits

verwendet wird;

eine erste Mehrzahl der Zeitfenster reserviert sind, um nur Datenbits von dem ersten isochronen Datenstrom zu übertragen, und zweite vorgegebene Zeitfenster verwendet werden, um Datenbits von dem zweiten Datenstrom zu übertragen; und Zuordnen der Daten der ersten isochronen Datenquelle/Senke zu der ersten vorgegebenen Mehrzahl Zeitfenster, um eine vorgegebene Datenbandbreite für die isochronen Daten von der ersten isochronen Datenquelle/Senke vorzusehen;

gekennzeichnet durch

Detektion einer Differenz zwischen einer tatsächlichen Rahmenlänge, die durch einen lokalen Oszillator eingestellt wird, und einer Rahmenlänge, die durch ein Netzwerkbezugstaktsignal bestimmt wird, und durch Kürzen oder Verlängern eines Rahmens durch Entfernen oder Hinzufügen eines Zeitfensters in Übereinstimmung damit, ob der lokale Oszillator relativ zum Bezugstakt langsam oder schnell ist.

Revendications

1. Appareil destiné à émettre de l'information vers un support physique (46) d'un réseau local comprenant: un moyen (50) destiné à recevoir une pluralité d'entrées, incluant au moins un premier flux de données isochrones et un second flux de données ; et un moyen (78) destiné à multiplexer dans le temps l'information provenant de ladite pluralité d'entrées en émettant des trames temporelles suivant une combinaison se répétant sensiblement sans pause entre lesdites trames temporelles, chaque trame temporelle comportant une pluralité de tranches de temps, les trames temporelles ayant des durées sensiblement égales, et chaque tranche de temps étant utilisée pour émettre un nombre prédéterminé de bits de données ; dans lequel des premières, prédéterminées, desdites tranches de temps de chaque trame temporelle sont utilisées pour émettre des bits de données issus dudit premier flux de données isochrones et des secondes, prédéterminées, desdites tranches de temps de chaque trame temporelle sont utilisées pour émettre des bits de données issus dudit second flux de données ; caractérisé par un moyen destiné à détecter une différence entre une longueur de trame réelle fixée par un oscillateur local et une longueur de trame déterminée par un signal d'horloge de référence de réseau et à raccourcir ou allonger une trame en supprimant ou ajoutant une tranche de temps selon que l'oscillateur local est lent ou rapide par rapport au signal d'horloge de référence.
2. Appareil selon la revendication 1, dans lequel ledit second flux de données est un flux de données à source non isochrone, dans lequel ledit premier flux de données isochrones est émis à une première cadence de données et dans lequel ladite première cadence de données est indépendante d'interruptions dans le transfert dudit second flux de données.
3. Appareil selon la revendication 1, dans lequel ledit second flux de données est constitué de données ayant pour source un réseau de type Ethernet ayant un certain potentiel d'événements de collision de données, et dans lequel ledit premier flux de données isochrones est émis à une cadence de données qui est sensiblement inchangée en réponse à l'occurrence d'un événement de collision de données dans ledit second flux de données.
4. Appareil selon la revendication 1, dans lequel ledit second flux de données est constitué de données ayant pour source un anneau à jetons, ayant un certain potentiel d'événements de perte de jeton, et dans lequel ledit premier flux de données isochrones est émis à une cadence de données qui est sensiblement inchangée en réponse à l'occurrence d'un événement de perte de jeton dans ledit second flux de données.
5. Appareil selon la revendication 1, dans lequel ledit second flux de données est un flux de données à source non isochrone, et dans lequel au moins une desdites tranches de temps est utilisée pour émettre des données issues dudit flux de données non isochrones pendant au moins une première pluralité de trames temporelles, et ladite une desdites tranches de temps est utilisée pour émettre un signal sans données pendant une seconde pluralité de trames temporelles.
6. Appareil selon la revendication 5, dans lequel ledit moyen destiné au multiplexage temporel émet une combinaison prédéterminée de ladite première pluralité de trames et de ladite seconde pluralité de trames.
7. Appareil selon la revendication 6, dans lequel ladite combinaison prédéterminée est une alternance entre ladite première pluralité de trames et ladite seconde pluralité de trames.

8. Appareil selon la revendication 1, dans lequel ledit second flux de données est un flux de données à source non isochrone, et dans lequel, dans une même trame temporelle, il n'y a pas plus de deux tranches de temps entre deux quelconques tranches de temps séquentielles à source non isochrone.

9. Appareil selon la revendication 1, dans lequel, dans une même trame temporelle, il n'y a pas plus de cinq tranches de temps entre deux quelconques tranches de temps remplies à source isochrone.

10. Appareil selon la revendication 1, dans lequel les première et deuxième tranches de temps dans chacune desdites pluralités de trames sont utilisées pour émettre des symboles de synchronisation de trame.

11. Appareil selon l'une quelconque des revendications précédentes et comprenant une pluralité de noeuds, chacun connecté à un noyau, et comprenant en outre un moyen destiné à recevoir un signal d'horloge de référence, un moyen destiné à propager ledit signal d'horloge de référence jusqu'audit noyau et jusqu'à chacun desdits noeuds, et un moyen destiné à initialiser l'émission de ladite trame temporelle à un instant prédéterminé en relation avec ledit signal d'horloge de référence propagé.

12. Appareil selon la revendication 11, et incluant un moyen destiné à recevoir un second signal d'horloge de référence, et un moyen destiné à choisir ledit premier signal d'horloge comme signal pour utilisation par ledit moyen destiné à initialiser l'émission.

13. Appareil selon la revendication 12, comprenant en outre une mémoire tampon destinée à mémoriser des données émises sur ledit support physique pour permettre la sortie, dudit réseau, synchronisée avec ledit second signal d'horloge de référence.

14. Procédé destiné à émettre de l'information vers un support physique couplé à une pluralité de sources et/ou de collecteurs de données d'un réseau local, incluant au moins une première source ou collecteur de données isochrones, le procédé comprenant :

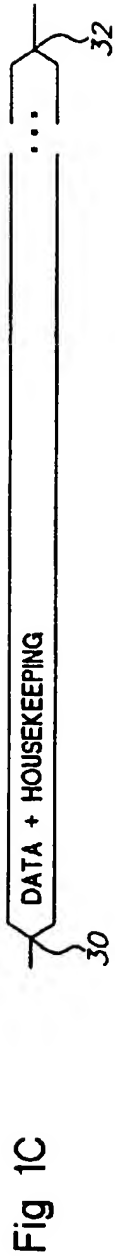
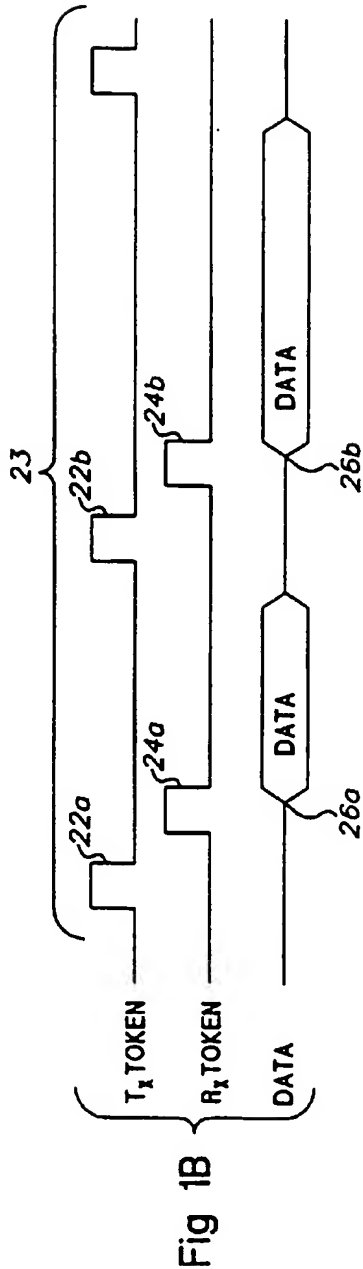
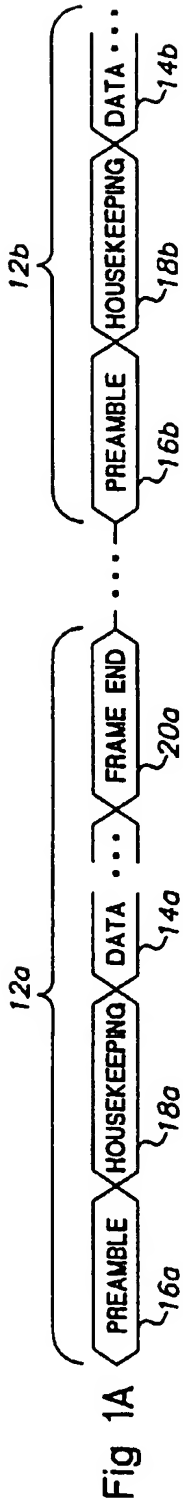
la réception, à une première de ladite pluralité de sources et/ou collecteurs de données, d'une pluralité d'entrées, incluant au moins un premier flux de données isochrones et un second flux de données, ledit premier flux de données isochrones incluant au moins des données issues de ladite première source ou collecteur de données isochrones ;

le multiplexage dans le temps d'information provenant de ladite pluralité d'entrées en émettant ladite information vers ledit support physique en une combinaison se répétant de trames temporelles sensiblement sans pause entre lesdites trames temporelles, chaque trame temporelle comportant une pluralité de tranches de temps, chaque trame temporelle ayant une durée sensiblement égale, et chaque tranche de temps étant utilisée pour émettre un nombre prédéterminé de bits de données ;

une première pluralité desdites tranches temporelles étant réservées pour émission seulement de bits de données issus dudit premier flux de données isochrones et des secondes, prédéterminées, desdites tranches de temps étant utilisées pour émettre des bits de données issus dudit second flux de données ;

et l'affectation de données issues de ladite première source ou collecteur de données isochrones, à certaines prédéterminées de ladite première pluralité de tranches de temps pour fournir une largeur de bande de données prédéterminée pour des données isochrones issues de ladite première source ou collecteur de données isochrones ;

caractérisé par la détection d'une différence entre une longueur de trame réelle fixée par un oscillateur local et une longueur de trame déterminée par un signal d'horloge de référence de réseau et par le raccourcissement ou l'allongement d'une trame en supprimant ou ajoutant une tranche de temps selon que l'oscillateur local est lent ou rapide par rapport au signal d'horloge de référence.



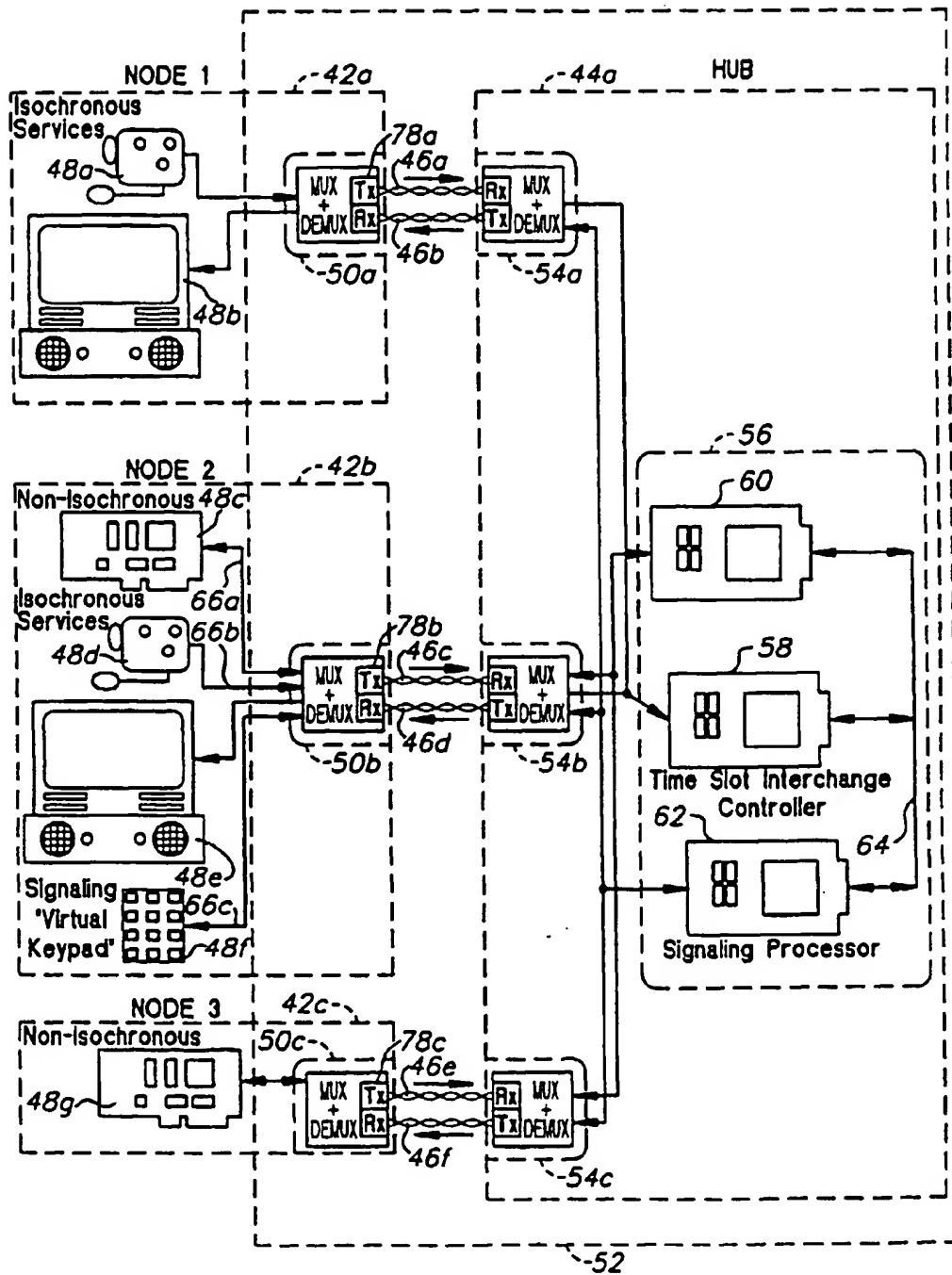


Fig 2

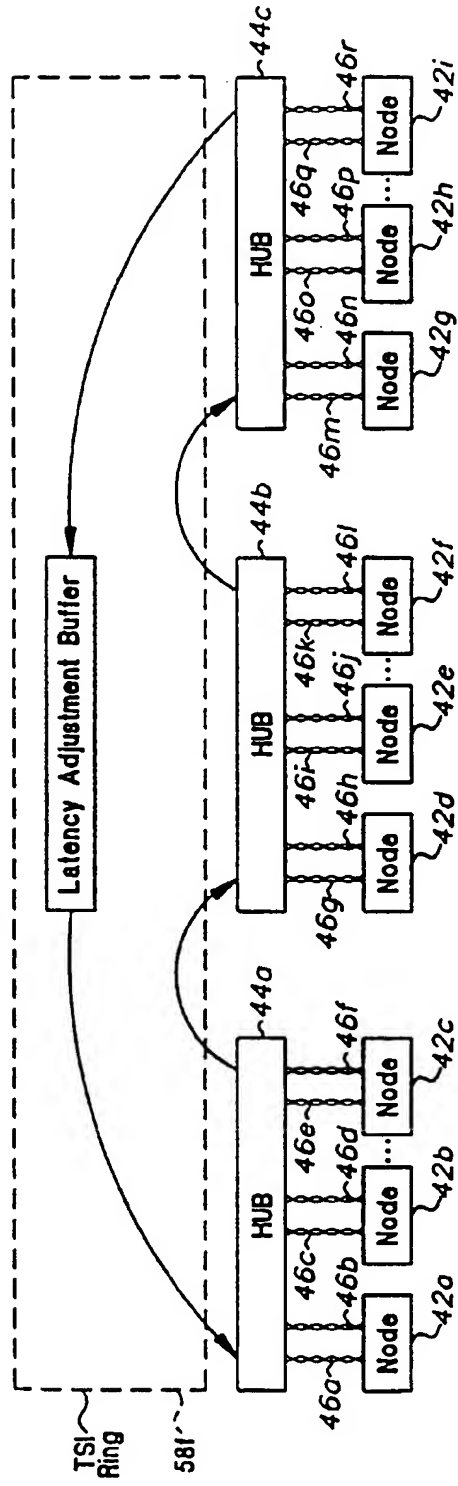


Fig 3

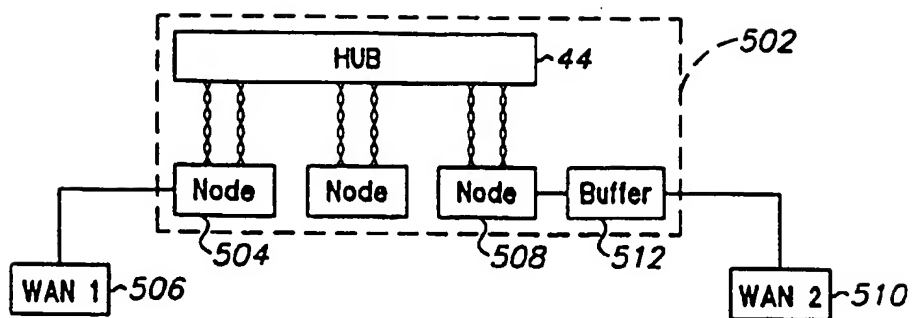


Fig 4

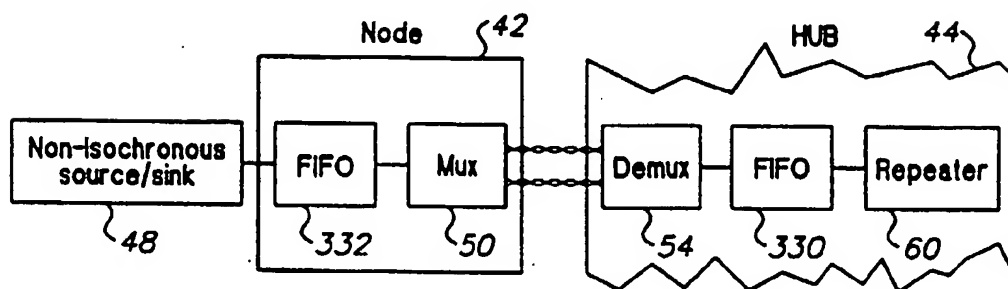


Fig 5

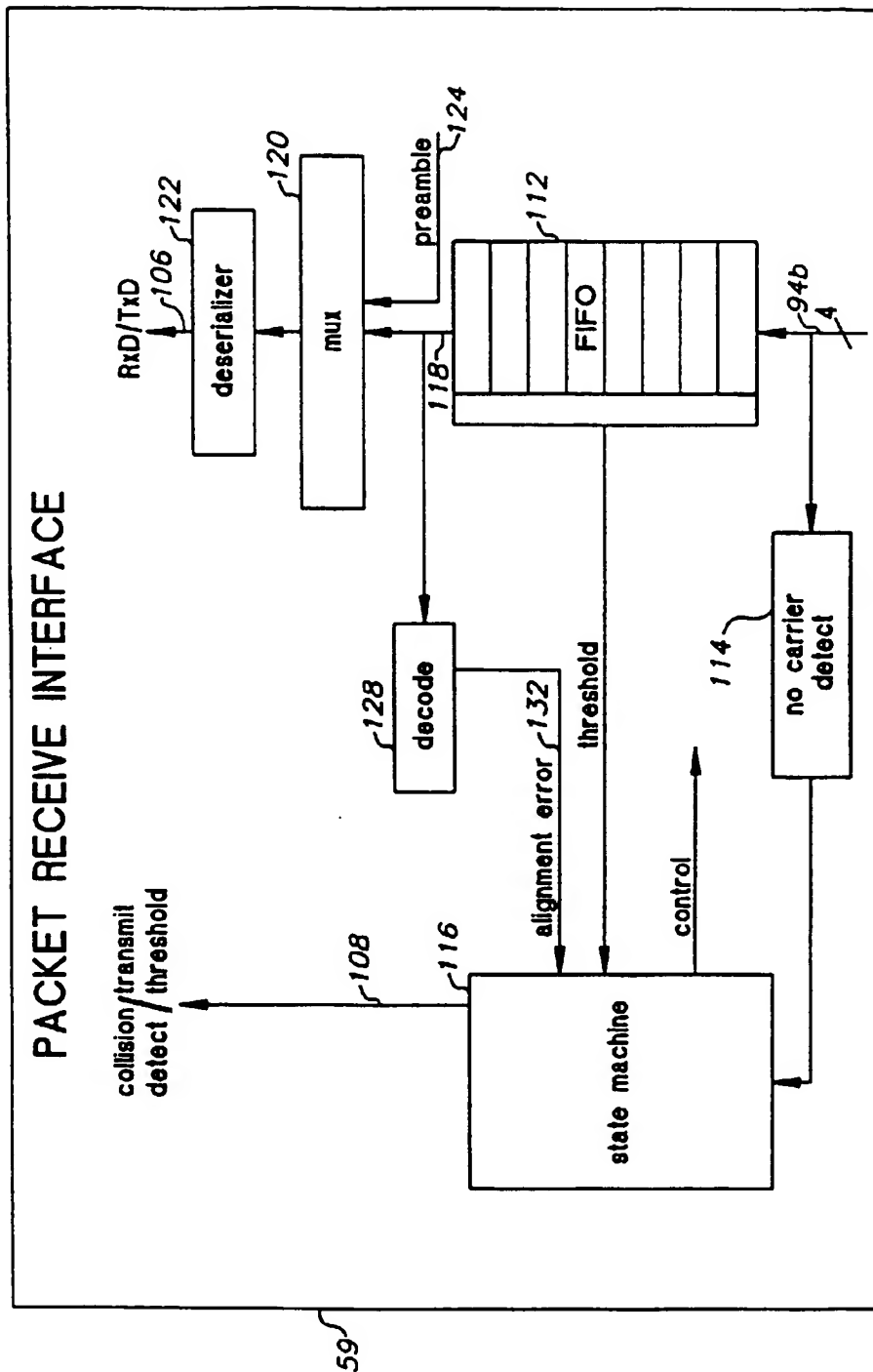


Fig 6

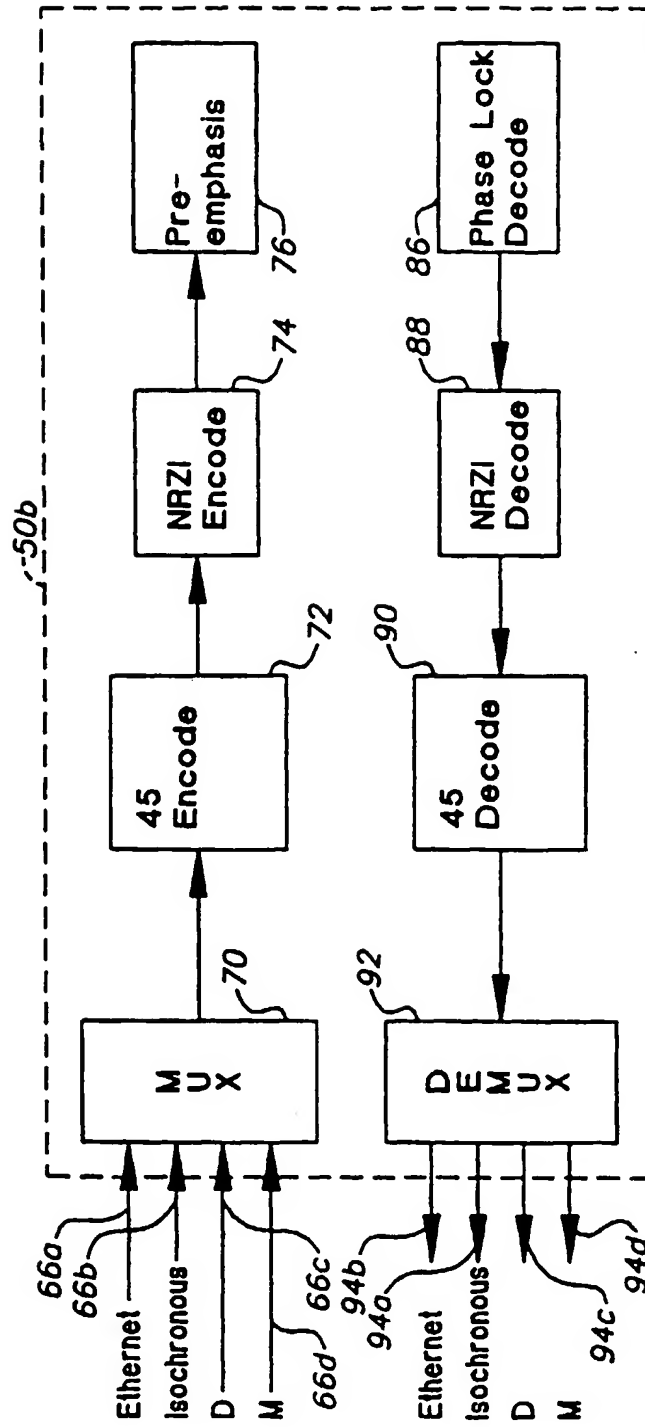


Fig 7

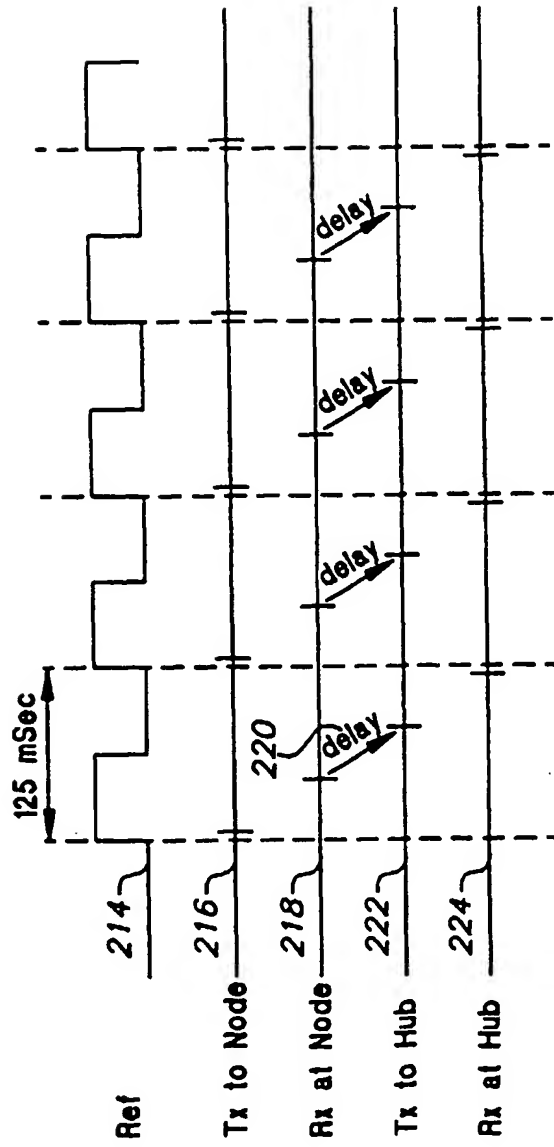


Fig 8